

# Search

Finding Solutions in Life since Brutus

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Computer Science: Artificial Intelligence

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Mathematical Proof

Search

Uninformed Search

Informed Search

# Proof

- Series of steps that logically follow from step to conclusion
- Showing validity
- To disprove, just show counterexample

# Enumeration

- “Brute Force”
- Just list every possibility
- Obviously only applies to finite, enumerable sets
- Useful for certain applications, like logic (truth tables)

# Direct

- Start with premises, and get to goal
- For equivalence, remember to prove both ways

# Proof by Contradiction

- Prove that the opposite of our statement is false
- Assume goal is false
- Prove a contradiction
- We essentially used this when doing automated logical resolution
- Examples
  - Infinitude of Primes
  - $\sqrt{2}$  is irrational

# Induction

- “Recursion in math”
- Requires partial ordering (ordered, and bounded at one end)
- Example
  - $0+1+\dots+n = n(n+1)/2$

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# Search

- High-level definition: Examining future options to determine what immediate action should be taken
- Assumptions
  - Environment Observable (not hidden)
  - Environment known
  - Deterministic actions

# Problem Formulations

- States
- Initial State
- Actions
- Transition Model
- Goal Test
- Path Cost
  - Step Cost

# Search

- Better definition: Looking for a *sequence of actions* to take that achieves the *goal*. That sequence of actions is called the *solution*.
- Initial State, actions, transition model are, together, called the *state space*.
  - All reachable states from the initial state
- All state spaces can be represented as a graph

# Search Terminology

- Incremental (start with nothing, and add states)
- Complete-state (start with everything) and rearrange, or delete

# Search Tree

- All searches can be represented as a search tree of nodes
- Nodes are not states. Nodes represent a state in the search
- Expanding nodes
- Leaf nodes of search together, is the *frontier*
- Can keep *explored set* to not visit repeat.

# General Graph Search

- See Russell/Norvig figure 3.7
- Different searches all about how you choose next leaf to expand.
- Nodes need to store
  - Corresponding state
  - Parent
  - Action taken by parent to get here
  - Path Cost ( $g(n)$ )

# Breath-First Search (BFS)

- Choose next node by putting nodes in Queue
- Ignores all costs

# Depth-First Search (DFS)

- Replace the queue in BFS with a stack



# Uniform-Cost Search

- Also called Dijkstra's Algorithm
- Expand node with lowest  $g(n)$

# Iterative Deepening DFS

- Go one layer at a time, doing DFS every time
- Repeats nodes

# Bidirectional Search

- Start search from both start and goal.
- This works because if  $e$  is avg number of edges at a node,
  - $e^t$  nodes expanded in  $t$  steps.
  - If  $2 < e$ , then  $2(e^t) < e(e^t)$

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# Informed Search

- Uses a *heuristic function*,  $h(n)$ 
  - Intuition about the world
  - Estimated cost of cheapest path from the state at node  $n$  to goal
  - Let  $C=g+h$
  - $h^*$ =oracle heuristic
  - $C^*$  = Optimal cost

# Detour: Distance

- $d$  is a distance iff  $d: X, Y \rightarrow \mathbf{R}$ , s.t.
  - $d(p, q) \geq 0$
  - $d(p, q) = 0$  iff  $p = q$
  - $d(p, q) = d(q, p)$       Commutativity
  - $d(p, q) \leq d(p, r) + d(r, q)$       Triangle Inequality

# Detour: Distance

- Euclidean Distances
  - Square roots of sum of squares ( $L_2$  norm)
  - Manhattan distance ( $L_1$  norm)
  - $L_{\infty}$  norm (Chebyshev distance) =  $\max_i(\text{abs}(p_i - q_i))$
- Non-Euclidean
  - Hamming distance
  - Edit distance
  - Cosine distance
  - Jaccard distance



# Evaluation

- **Completeness:** Will check all nodes to find solution
- **Optimality:** The solution returned will be optimal
- **Optimally Efficient:** For another algorithm using same cost function and same heuristic, you cannot expand fewer nodes.

# Best-First Search

- Give every node an  $f(n)$  to calculate a priority.
- Pick the one with the best priority
- With uniform cost search, notice  $f(n)=g(n)$
- Greedy Best-first search
  - Intuition: Expand node closest to goal
  - Complete: Yes
  - Optimal: No

# A\* Search

- Set  $f(n) = g(n) + h(n)$
- Intuitively,  $f(n)$  is estimated cost from initial state to goal through the current node
- Conditions
  - $h(n)$  must be *admissible heuristic*
  - Consistency (monotonicity)
  - $h(n) \leq c(n, a, n') + h(n')$

$A^*$

- Optimal
- Complete
- Optimally Efficient
- However, can be costly

# Alternatives

- IDA\* (Iterative Deepening A\*)
- RFBS (Recursive BFS)
- SMA (Simple Memory-Bounded A\*)

# Search

- Domination
  - For any node  $n$ , if  $h_2 \geq h_1$  then  $h_2$  dominates  $h_1$ .
  - Since we can't pass  $C^*$ , dominant heuristics are good
  - A weak heuristic takes us closer to uniform-cost search
- Generating heuristics
  - Multiple heuristics: can take  $\max\{h_1, h_2, h_3, \dots\}$
  - Cost of optimal solution to a relaxed problem is an admissible heuristic for the original problem